A Study of Divide and Distribute Fixed Weights and its Variants

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Based on the PoS submission found at http://crcodel.com/research/ddfw_pos.pdf and the larger research thesis found at http://crcodel.com/research/ddfw_thesis.pdf.

We studied the Divide and Distribute Fixed Weights (DDFW) stochastic local search algorithm

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In local minima, DDFW distributes weight from satisfied to unsatisfied clauses

We studied how to best flip variables and distribute weight by testing DDFW against modern hard benchmarks

The DDFW algorithm

Variable flip variants

Weight redistribution variants

The DDFW algorithm

All clauses receive an initial weight ($w_{init}[=8]$)

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DDFW flips variables which most reduce the unsatsfied clause weight (W_U)

In local minima, DDFW moves weight from one satisfied neighbor to each unsatisfied clause

DDFW is the only SLS algorithm in the UBCSAT framework to efficiently solve the n = 7824 Pythagorean triples instance

The DDFW algorithm as pseudocode

Algorithm 1: DDFW

1	Input: CNF formula ${\cal F}$					
2	Set all clause weights to $w_{ m init}=8$					
3	$\alpha \leftarrow randomly generated truth assignment$					
4	4 for MAX-FLIPS times do					
5	if assignment α satisfies $\mathcal F$ then return α					
6	if flipping a variable reduces W _U then					
7	Flip a literal that reduces W_U the most					
8	else					
9	for each unsatisfied clause C _i do					
10	$C_k \leftarrow maximum-weight neighbor of C_i$					
11	if weight of $C_k > w_{\text{init}}$ then					
12	Transfer a weight of 2 from C_k to	C_i				
13	else	,				
14	Transfer a weight of 1 from C_k to	C_i				
15	15 return "No satisfying assignment"					

The DDFW algorithm as pseudocode

Algorithm 2: DDFW



Benchmark set

Ten encodings of matrix multiplication challenges

Ten random 3-SAT instances from the 2018 SAT Competition

Two encodings of the n = 7824 Pythagorean triples problem

Two encodings of asias and three of Steiner triple problems

All CNFs can be found at https://github.com/marijnheule/benchmarks and http://satcompetition.org The DDFW algorithm

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DDFW greedily flips variables which reduce W_U the most

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Idea: flip W_U -reducing variables probabilistically

We investigated a uniform and a weighted probability distribution

Variable flip variant experimental results



Averaged over all problem instances. The greedy (original) method performed significantly better than the variants.

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DDFW distributes 1 or 2 units of weight between clauses

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We can generalize to a linear rule:

Algorithm 4: Linear weight transfer rule

1 if weight of unsat neighbor $C_k > w_{init}$ then 2 | Transfer a weight of $a_> \times w(C_k) + c_>$ from C_k to C_j 3 else

4 | Transfer a weight of $a_{\leq} \times w(C_k) + c_{\leq}$ from C_k to C_j

Weight redistribution variants experimental results

 $w_{\rm init}=100.$ Results averaged over all instances and 100 runs per instance, five million flip timeout

Distribution policy	Avg lowest unsat	Solve %
Original DDFW	36.57	0.85
$(c_{>}, c_{\leq}) = (10, 25)$	29.16	1.7
$(a_{>},a_{\leq}) = (0.05, 0.05)$	23.82	2.96
$(a_{>},c_{>}) = (a_{\leq},c_{\leq}) = (0.1,5)$	22.05	2.67

Linear rule performed about 40% better

Weight redistribution variants experimental results

Parameter searches across the matrix multiplication challenges



Note greater effect of $a_>$ on shape of right plot, while c_\leq determines shape of left plot

Weight redistribution variants continued

Can also distribute weight from entire neighborhoods

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Two methods tested: apply linear rule to each clause or to clause proportional to clause weight

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Experimental results disappointing but show some promise

Distribution policy	Avg lowest unsat	Matrix lowest unsat
Original DDFW	36.57	57.02
Proportional	46.91	27.0
Direct	45.29	33.91

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Spreading weight across more clauses in a neighborhood could cause $\rm DDFW$ to escape local minima faster