# On Irrelevant Literals in Pseudo-Boolean Constraint Learning

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#### A linear pseudo-Boolean (PB) constraint is of the form

$$\sum_{j} a_{j} l_{j} \vartriangle k$$

#### where

- $\forall j, a_j \in \mathbb{Z}$
- $\forall j, l_j$  is a literal (i.e. a Boolean value)
- $\bullet \ {\bigtriangleup } \in \{<,\leqslant,=,\geqslant,>\}$
- $k \in \mathbb{Z}$  is the degree of the constraint

Example: 
$$3a - 2\overline{b} + c - 4d \leq -1$$

#### Normalized PB constraints

$$\sum_{j} a_{j} l_{j} \ge k \text{ with } \forall j, a_{j} \in \mathbb{N} \text{ and } k \in \mathbb{N}$$

Example:  $3a + 2\overline{b} + c \ge 3$ 

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Clauses

$$\sum\limits_{j} l_{j} \geqslant 1$$
  
Example:  $a + \overline{b} + c \geqslant 1$ 

# Generalized Resolution [Hooker, 1988]

Most PB solvers use the following rules to learn new constraints (a.k.a. no-goods) when they encounter a conflict, so as not to do the same mistake again

$$\frac{al + \sum_{i=1}^{n} a_i l_i \ge d_1 \qquad b\bar{l} + \sum_{i=1}^{n} b_i l_i \ge d_2}{\sum_{i=1}^{n} (ba_i + ab_i) l_i \ge bd_1 + ad_2 - ab}$$
(clashing addition)

$$\frac{\sum_{i=1}^{n} a_i l_i \ge d}{\sum_{i=1}^{n} \min(a_i, d) l_i \ge d}$$
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This proof system is (theoretically) more powerful than classical resolution: its proofs may be exponentially shorter

Consider the following constraints

$$\begin{split} \chi_1 &: \bar{a} + \bar{b} + f \geqslant 2\\ \chi_2 &: 3\bar{x} + a + b + d + e \geqslant 4\\ \chi_3 &: 4a + 2b + 2c + x \geqslant 5 \end{split}$$

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$$f = 0@1 \cdot 4$$

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$$\chi_{1} \qquad \chi_{2}$$

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This constraint is learned because it propagates a to 1 at level 0

 $13a + 7b + 6c + d + e \ge 16$ 

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 $13a + 7b + 6c \ge 14$ 

# Irrelevant Literals in Practice (in Sat4j)



- Number of irrelevant literals in Sat4j-CP's first 5,000 learned constraints
- Experiments conducted on the 777 decision benchmarks from PB'16
- Sat4j as an example of Generalized-Resolution-based solver

*RoundingSat* uses a different approach, which mainly consists in using the division rule instead of saturation

$$\frac{\sum_{i=1}^{n} a_i l_i \ge d \qquad \alpha > 0}{\sum_{i=1}^{n} \left\lceil \frac{a_i}{\alpha} \right\rceil l_i \ge \left\lceil \frac{d}{\alpha} \right\rceil}$$
(division)

Consider the following constraints:

 $\chi_1 : 2\bar{c} + 2\bar{d} + b + \bar{e} \ge 4$   $\chi_2 : 3a + 3b + c + d + e \ge 4$  $\chi_3 : 2\bar{a} + b + e \ge 2$ 

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$$e = 1@1 \cdot$$



$$\chi_1 : 2\bar{c} + 2\bar{d} + \mathbf{b} + \bar{\mathbf{e}} \ge 4$$
  

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We have falsified  $\chi_3!$  Before applying clashing addition,  $\chi_2$  is weakened on *e* and divided by 3

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$$a+b+c+d \ge 1$$

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Observe how c and d become irrelevant, and then relevant again, and how they prevent the inference of the stronger constraint  $a + b \ge 1$ 

# Irrelevant Literals in Practice (in RoundingSat)



- Number of irrelevant literals in RoudingSat's first 100,000 weakened constraints
- Experiments conducted on the 777 decision benchmarks from PB'16

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  - $\equiv$  15*a* + 10*b* + 10*c*  $\geq$  15

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 $\equiv$ 

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*Efficient data structures implemented in PB solvers cannot be used when cardinality constraints are hidden* 

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So, in practice, performing a complete removal on all inferred constraints seems out of reach

But we can still consider an incomplete approach, treating only small learned constraints, and optimizing the detection phase

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This test can be achieved by verifying that this formula is unsatisfiable

$$\chi \wedge \neg \Big( 13a + 7b + 6c + d \ge 16 \Big)$$

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The relevance of a literal can be checked using a PB solver

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13a + 7b + 6c + d = (16 - 1) = 15

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Its ability to compute "efficiently" all possible sums is crucial here

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If a literal is relevant, so it is for all literals with a greater coefficient

The previous observations lead to the following algorithm

Algorithm 1: detect-and-remove-irrelevant-literals

**Input:** A non-valid pseudo-Boolean constraint  $\chi$ 

**Output:** The constraint  $\chi$ , in which all irrelevant literals are removed

 $\begin{array}{c|c} \textbf{foreach coefficient } c \ of \ the \ constraint \ in \ ascending \ order \ \textbf{do} \\ \hline \\ Choose a \ literal \ l \ having \ coefficient \ c \\ \hline \\ \textbf{if } \ depends(l,\chi) \ \textbf{then} \\ | \ \ \textbf{return } \chi \\ \hline \\ \textbf{end} \\ \hline \\ Remove \ all \ literals \ having \ coefficient \ c \ from \ \chi \\ \hline \\ Update \ the \ degree \ of \ \chi \\ \hline \\ Saturate \ \chi \\ \textbf{end} \\ \end{array}$ 

#### **Experiments**

#### Settings

- Quadcore bi-processors Intel XEON E5-5637 v4 (3.5 GHz)
- 128 GB of memory
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#### **Experimented approaches**

- Detection and removal implemented in Sat4j-CuttingPlanes
- Using a dynamic programming algorithm or a PB solver (5-second timeout per call)
- Only applied to learned constraints having less than 1,000 literals and a degree less than 20,000

# **Experimental Results: Detection using Dynamic Programming**



- Instances not solved at all are not presented
- $\bullet$  A  $\,\circ\,$  stands for SATISFIABLE and a  $\,\square\,$  stands for UNSATISFIABLE

### Experimental Results: Detection with Sat4j-CP



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### **Experimental Results: Detection with Sat4j-Res**



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## **Experimental Results: Detection Runtime**



### Conclusion

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#### Perspectives

- Improve the detection algorithm to remove more literals
- Find a proof system guaranteeing not to infer such literals

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